

**TMA947/MMG621  
NONLINEAR OPTIMISATION**

**Date:** 15-04-14

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**Question 1**

(the simplex method)

- (2p) a) We first rewrite the problem on standard form. We introduce slack variables  $s_1$  and  $s_2$ . Consider the following linear program:

$$\begin{aligned} \text{minimize} \quad & z = -2x_1 - x_2 \\ \text{subject to} \quad & -x_1 + x_2 + s_1 = 1, \\ & x_1 - 2x_2 + s_2 = 2, \\ & x_1, x_2, s_1, s_2 \geq 0. \end{aligned}$$

In phase I the starting basis is  $(s_1, s_2)^T$ . Calculating the reduced costs for the non-basic variables  $x_1, x_2$  we obtain  $\tilde{\mathbf{c}}_N = (-2, -1)^T$ , meaning that  $x_1$  enters the basis. From the minimum ratio test, we get that  $s_2$  leaves the basis.

Updating the basis we now have  $(s_1, x_1)^T$  in the basis. Calculating the reduced costs, we obtain  $\tilde{\mathbf{c}}_N = (2, -5)^T$ . meaning that  $x_2$  enters the basis. From the minimum ratio test we get  $\mathbf{B}^{-1}\mathbf{N}_2 = (-1, -2)^T < \mathbf{0}$ , meaning that the problem is unbounded.

- (1p) b) A direction of unboundness is  $\mathbf{l}(\mu) = (2, 0, 3, 0)^T + \mu(2, 1, 1, 0)^T, \mu \geq 0$ .

**(3p) Question 2**

(consistency of linear systems)

Consider to linear program to

$$\begin{aligned} \text{minimize} \quad & -\mathbf{c}^T \mathbf{x}, \\ \text{subject to} \quad & \mathbf{A}\mathbf{x} \leq \mathbf{b}, \end{aligned} \tag{1}$$

and its standard form equivalent

$$\begin{aligned} \text{minimize} \quad & -\mathbf{c}^T \mathbf{x}^+ + \mathbf{c}^T \mathbf{x}^-, \\ \text{subject to} \quad & \mathbf{A}\mathbf{x}^+ - \mathbf{A}\mathbf{x}^- + \mathbf{s} = \mathbf{b}, \\ & \mathbf{x}^+ \geq \mathbf{0}, \mathbf{x}^- \geq \mathbf{0}, \mathbf{s} \geq \mathbf{0}. \end{aligned} \tag{2}$$

The dual of (2) is to

$$\begin{aligned}
 & \underset{\mathbf{p}}{\text{maximize}} && \mathbf{b}^T \mathbf{p}, \\
 & \text{subject to} && \mathbf{A}^T \mathbf{p} \leq -\mathbf{c}, \\
 & && -\mathbf{A}^T \mathbf{p} \leq \mathbf{c}, \\
 & && \mathbf{p} \leq \mathbf{0},
 \end{aligned} \tag{3}$$

and (3) is equivalent to

$$\begin{aligned}
 & \underset{\mathbf{y}}{\text{maximize}} && -\mathbf{b}^T \mathbf{y}, \\
 & \text{subject to} && \mathbf{A}^T \mathbf{y} = \mathbf{c}, \\
 & && \mathbf{y} \geq \mathbf{0}.
 \end{aligned} \tag{4}$$

If statement (a) holds, then the optimal objective value of (1) and (2) are bounded from below by  $-d$ . Hence, by there exists an optimal solution to the dual of (2), which is (3). Consequently, by strong duality (cf. Theorem 10.6) the optimal objective values of (3) and (4) are equal to that of (2), which is bounded from below by  $-d$ . This implies that, for (4), there exists a vector  $\mathbf{y} \geq \mathbf{0}$  such that  $\mathbf{A}^T \mathbf{y} = \mathbf{c}$  and  $-\mathbf{b}^T \mathbf{y} \geq -d$  (i.e.,  $\mathbf{b}^T \mathbf{y} \leq d$ ). This statement is the same as (b).

Conversely, if (b) holds then (3) has at least one feasible solution with objective value bounded from below by  $-d$ . Hence, by weak duality (cf. Theorem 10.4) every  $\mathbf{x}$  feasible in (1) (i.e.,  $\mathbf{Ax} \leq \mathbf{b}$ ) must satisfy  $-\mathbf{c}^T \mathbf{x} \geq -d$ . This implies statement (a).

### (3p) Question 3

(global optimality conditions)

This is Theorem 6.8.

### (3p) Question 4

(modelling)

The decision variables are:

$x_A$  = number of units of product A produced

$x_B$  = number of units of product B produced

$y_1 = 1$ , if additional time is used, 0 otherwise

$y_2 = 1$ , if more than ten units of product A is produced, 0 otherwise

Based on these definitions, the model is as follows:

$$\begin{aligned} & \text{maximize} && 200x_A + 400x_B - 1200y_1 \\ & \text{subject to} && 2x_A + 3x_B \leq 40 + 8y_1, \\ & && 100y_2 \geq x_A - 10, \\ & && x_B \geq 5y_2, \\ & && x_A, x_B \geq 0, \text{ integer} \\ & && y_1, y_2 \in \{0, 1\} \end{aligned}$$

The program is linear with integer variables.

### Question 5

(true or false)

- (1p) a) False. The directional derivative must be non-negative in *all* directions  $\mathbf{p}$ .
- (1p) b) False. The problem is feasible but may have an unbounded solution.
- (1p) c) True. This is a consequence of Theorem 4.23.

### (3p) Question 6

(nonlinear programming)

Letting  $\mu$  denote the Lagrange multiplier for the equality constraint, and  $\boldsymbol{\lambda} \in \mathbb{R}_+^n$  denote the vector of multipliers for the sign constraints, we obtain the Lagrangian

$$L(\mathbf{x}, \mu, \boldsymbol{\lambda}) := f(\mathbf{x}) + \mu \left( \sum_{j=1}^n x_j - r \right) - \boldsymbol{\lambda}^T \mathbf{x}.$$

Consider the optimality condition for  $x_j$ :

$$\frac{\partial L(\mathbf{x}, \mu, \boldsymbol{\lambda})}{\partial x_j} = \frac{\partial f(\mathbf{x})}{\partial x_j} - \mu - \lambda_j = 0, \quad j = 1, \dots, n.$$

Further, we have that  $\lambda_j^* x_j^* = 0$ , by complementarity. If  $x_j^* > 0$  then  $\lambda_j^* = 0$ , and hence  $\frac{\partial f(\mathbf{x}^*)}{\partial x_j} = \mu^*$  (hence a common partial derivative for *all* positive variables), while if  $x_j^* = 0$  then  $\frac{\partial f(\mathbf{x}^*)}{\partial x_j} = \mu^* + \lambda_j^*$ , which may be larger.

### (3p) Question 7

(gradient projection algorithm)

Denote the objective function by  $f(x_1, x_2) := 3x_1^2 - 2x_1x_2 + 2x_2^2$ , and the (box) feasible set by  $X$ . Then,  $\nabla f(\mathbf{x}) = (6x_1 - 2x_2, -2x_1 + 4x_2)^T$ . At the initial point  $\mathbf{x}^0 = (1, -2)^T$ , the gradient is  $\nabla f(\mathbf{x}^0) = (10, -10)^T$ . To determine step length  $\alpha^0$ , we apply the Armijo criterion supplied. We first try  $\alpha^0 = \bar{\alpha} = 1$  (as  $\beta^0 = 1$ ). Note that

$$\begin{aligned} \text{Proj}_X[\mathbf{x}^0 - \nabla f(\mathbf{x}^0)] &= \text{Proj}_X \left[ \begin{pmatrix} 1 - 10 \\ -2 + 10 \end{pmatrix} \right] = \begin{pmatrix} 0 \\ -1 \end{pmatrix} \\ \text{Proj}_X[\mathbf{x}^0 - \nabla f(\mathbf{x}^0)] - \mathbf{x}^0 &= \begin{pmatrix} 0 \\ -1 \end{pmatrix} - \begin{pmatrix} 1 \\ -2 \end{pmatrix} = \begin{pmatrix} -1 \\ 1 \end{pmatrix} \\ f(\text{Proj}_X[\mathbf{x}^k - \bar{\alpha}\beta^i \nabla f(\mathbf{x}^k)]) &= 3 \cdot 0 + 2 \cdot 0 - 2 \cdot (-1)^2 = 2 \\ f(\mathbf{x}^0) &= 3 \cdot 1^2 - 2 \cdot 1 \cdot (-2) + 2 \cdot (-2)^2 = 15 \\ \nabla f(\mathbf{x}^0)^T (\text{Proj}_X[\mathbf{x}^0 - \nabla f(\mathbf{x}^0)] - \mathbf{x}^0) &= (10 \quad -10) \begin{pmatrix} -1 \\ 1 \end{pmatrix} = -20. \end{aligned}$$

Hence, the Armijo criterion is satisfied, as  $2 \leq 15 + 0.2 \cdot (-20) = 11$ . Thus,  $\alpha^0 = \bar{\alpha} = 1$ , and iterate  $\mathbf{x}^1 = \text{Proj}_X[\mathbf{x}^0 - \nabla f(\mathbf{x}^0)] = (0, -1)^T$ .

For the next iteration, we have  $\nabla f(\mathbf{x}^1) = (2, -4)^T$ . Hence,

$$\mathbf{x}^1 - \alpha \nabla f(\mathbf{x}^1) = \begin{pmatrix} 0 \\ -1 \end{pmatrix} + \begin{pmatrix} -2\alpha \\ 4\alpha \end{pmatrix} = \begin{pmatrix} -2\alpha \\ -1 + 4\alpha \end{pmatrix}.$$

As a result,

$$\text{Proj}_X[\mathbf{x}^1 - \alpha \nabla f(\mathbf{x}^1)] = \begin{pmatrix} \max\{0, -2\alpha\} \\ \min\{-1, -1 + 4\alpha\} \end{pmatrix} = \begin{pmatrix} 0 \\ -1 \end{pmatrix} = \mathbf{x}^1, \quad \forall \alpha > 0,$$

and hence  $\mathbf{x}^1$  is a stationary point (KKT point). The gradient projection algorithm terminates because the termination criterion is met. Notice that the fact that  $\mathbf{x}^1$  is a stationary point can also be understood graphically, as  $-\nabla f(\mathbf{x}^1)$  lies in the cone generated by the normal vectors of the two active constraints ( $x_1 \geq 0$  and  $x_2 \leq -1$ ).

Finally, since  $f$  is convex (which can be verified by computing the Hessian) and  $X$  is convex, the stationary point  $\boldsymbol{x}^1$  is also optimal. This can be established via Theorem 4.23 together with (4.18), or Theorem 5.49 in the text.

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