

LECTURE 13: FEASIBLE DIRECTION METHODS

METHODS:

- UNCONSTRAINED METHODS
 - STEEPEST DESCENT
 - NEWTON'S METHOD
- METHODS FOR LPs
 - SIMPLEX
- METHODS FOR CONVEX PROBLEMS
 - SUBGRADIENT METHOD
- METHODS FOR ILP
- METHODS FOR CONSTRAINED PROBLEMS:

1. FEASIBLE DIRECTION METHODS (TODAY)
2. PENALTY METHODS (NEXT LECT.)

CONSIDER THE PROBLEM $f^* = \inf f(x)$
S.T. $x \in \underline{X}$

WHERE \underline{X} IS NONEMPTY, CLOSED & CONVEX, AND $f \in C^1$ ON \underline{X}

- IDEA: GENERALIZE UNCONSTRAINED OPTIMIZATION METHODS

GENERAL FEASIBLE-DIRECTION DESCENT METHOD

0: INITIALIZE $x^0 \in \bar{X}$, SET $k := 0$

1: FIND A FEASIBLE DESCENT DIRECTION $p_k \in \mathbb{R}^n$ SUCH THAT THERE EXISTS $\bar{\alpha} > 0$:

$$x_k + \alpha p_k \in \bar{X}, \quad \forall \alpha \in [0, \bar{\alpha}]$$

$$f(x_k + \alpha p_k) < f(x_k), \quad \forall \alpha \in (0, \bar{\alpha}]$$

2: DETERMINE STEP LENGTH $\alpha_k > 0$: $x_k + \alpha_k p_k \in \bar{X}$

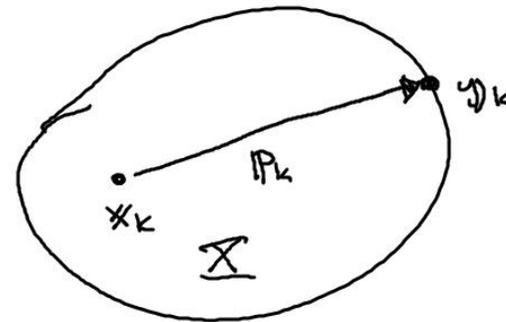
$$f(x_k + \alpha_k p_k) < f(x_k)$$

3: UPDATE $x_{k+1} = x_k + \alpha_k p_k$

4: CHECK TERMINATION CRITERIA. OTHERWISE GOTO 1 AND LET $k = k+1$

- JUST AS LOCAL AS UNCONSTRAINED METHODS
- SEARCH DIRECTIONS ARE OFTEN ON THE FORM

$$p_k = y_k - x_k$$
 WHERE $y_k \in X$ SOLVES SOME OTHER PROBLEM
- LINE SEARCH IS SIMILAR TO UNCONSTRAINED CASE. JUST DONT GO TO FAR!
- TERMINATION CRITERIA ARE OFTEN BASED ON KKT CONDITIONS



- FOR GENERAL \mathcal{X} , FINDING FEASIBLE DESCENT DIRECTIONS CAN BE TRICKY. THE SAME HOLDS FOR STEP LENGTHS.
- IF \mathcal{X} IS POLYHEDRON \Rightarrow SEARCH DIRECTIONS AND STEP LENGTHS ARE EASY TO FIND
- IF \mathcal{X} IS POLYHEDRON \Rightarrow KKT CONDITIONS NECESSARY.
 x^* LOCAL MIN $\Rightarrow x^*$ KKT POINT

THE FRANK-WOLFE METHOD

- ASSUME $\min_{x \in \mathcal{X}} f(x)$ WHERE \mathcal{X} IS POLYHEDRON.
S.T. $x \in \mathcal{X}$ AND $f \in C^1$ ON \mathcal{X}

- OPTIMALITY CONDITIONS:

x^* LOCAL MIN OF f OVER $\mathcal{X} \Rightarrow x^*$ STATIONARY

$$\Rightarrow \min_{x \in \mathcal{X}} \nabla f(x^*)^T (x - x^*) = 0$$

- IF $\min_{x \in \mathcal{X}} \nabla f(x^*)^T (x - x^*) \neq 0$ $\Rightarrow x^*$ IS NOT LOCAL MIN

\Rightarrow CAN CONSTRUCT FEASIBLE DESCENT DIRECTION AT x^*

FRANK-WOLFE

- AT $x_k \in \mathcal{X}$, IF

$$\left\{ \begin{array}{l} \min_{x \in \mathcal{X}} \nabla f(x_k)^T (x - x_k) < 0 \\ y_k \in \arg \min_{x \in \mathcal{X}} \nabla f(x_k)^T (x - x_k) \end{array} \right.$$

LINEAR PROGRAM!

- \mathcal{X} POLYHEDRON

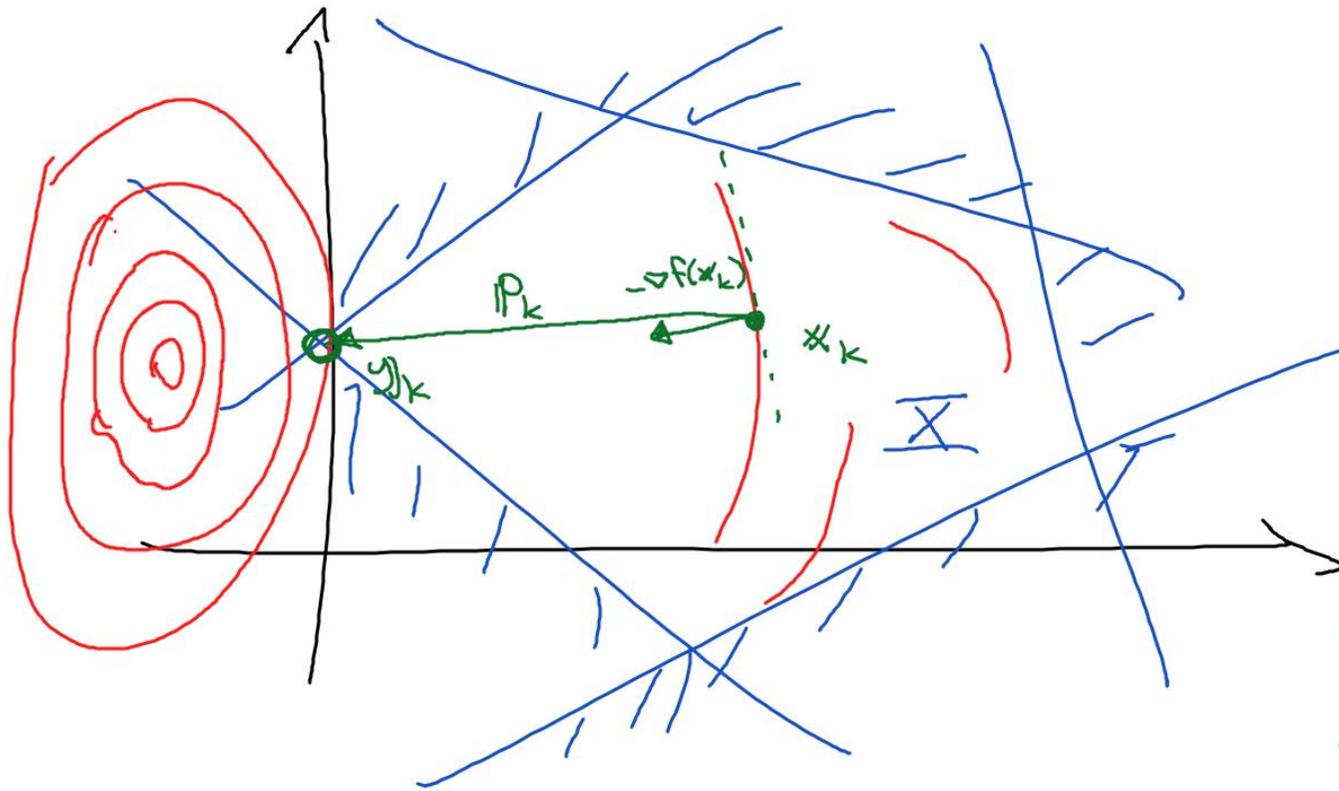
- OBJECTIVE LINEAR W.R.T. x

- EASY TO SOLVE

- THEN $p_k = y_k - x_k$ WILL BE A FEASIBLE DESCENT DIRECTION

- y_k WILL BE AN EXTREME POINT!

- SEARCH DIRECTION WILL ALWAYS BE TOWARDS AN EXTREME POINT
IN THE FRANK - WOLFE METHOD



- AT ITERATION x_k
WE COMPUTE $\nabla f(x_k)$.

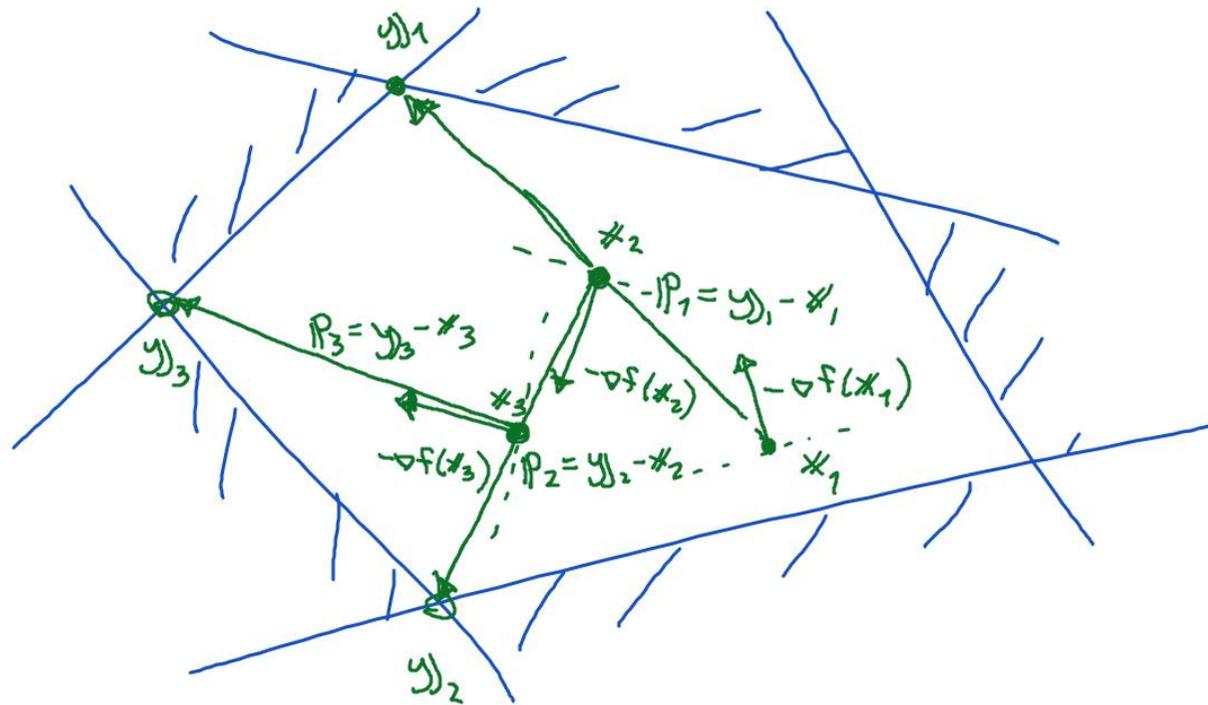
- NOW SOLVE THE LP

$$\min \underbrace{\nabla f(x_k)^T}_{=c} x$$
 s.t. $x \in X$

- LET y_k BE SOLUTION

- LET SEARCH DIRECTION
BE

$$p_k = y_k - x_k$$



FRANK-WOLFE METHOD

0: LET $x_0 \in X$. SET $k := 0$

1: FIND AN OPTIMAL SOLUTION TO THE PROBLEM (LP)

$$\min_{x \in X} z_k(x) = \nabla f(x_k)^T (x - x_k)$$

$$\text{LET } p_k = y_k - x_k$$

2: LINE SEARCH . FIND $\alpha_k \in [0, 1]$

3: $x_{k+1} = x_k + \alpha_k p_k$

4: IF $z_k(y_k)$ IS CLOSE TO 0 OR α_k IS CLOSE TO 0
WE TERMINATE , OTHERWISE LET $k = k+1$ AND GOTO 1.

WE WILL NEVER BECOME
INFEASIBLE



CONVERGENCE:

- SUPPOSE \bar{X} IS NONEMPTY POLYTOPE AND $f \in C^1$ ON \bar{X} .
IF x_k IS CHOSEN ACCORDING TO ARMIJOS RULE THEN
 $\{x_k\}$ IS BOUNDED AND EVERY LIMIT POINT IS STATIONARY!
- IF f IS CONVEX ON \bar{X} , THEN EVERY LIMIT POINT IS GLOBAL OPTIMUM.

IF f IS A CONVEX FUNCTION

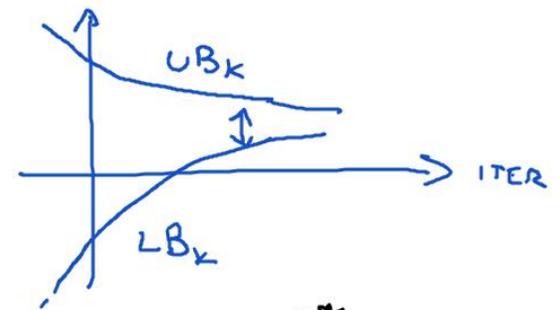
FOR EACH k , AND ALL $x \in \bar{X}$ IT HOLDS THAT

$$f(x) \geq f(x_k) + \nabla f(x_k)^T (x - x_k)$$

[SINCE f IS CONVEX]

$$\geq \underbrace{f(x_k) + \nabla f(x_k)^T (y_k - x_k)}$$

GET THIS IN EACH ITERATION



$$\Rightarrow f^* \geq f(x_k) + \nabla f(x_k)^T (y_k - x_k)$$

- SO IN EACH ITERATION WE GET A LOWER BOUND ON f^*

- KEEP BEST BOUND:

$$UB_{k+1} = \min [UB_k, f(x_k)] \quad [= f(x_k)]$$

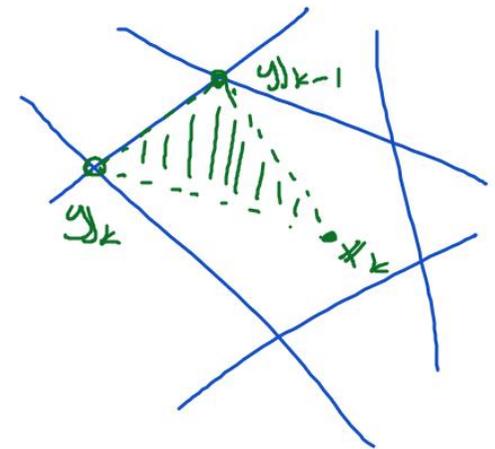
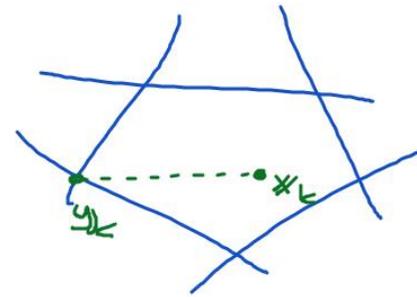
$$LB_{k+1} = \max [LB_k, f(x_k) + \nabla f(x_k)^T (y_k - x_k)]$$

SIMPLICIAL DECOMPOSITION

- FRANK-WOLFE METHOD FINDS NEXT ITERATE BETWEEN CURRENT ITERATE AND ONE EXTREME POINT

- IN EACH ITERATION A NEW EXTREME POINT IS COMPUTED.

- SIMPLICIAL DECOMPOSITION METHOD REMEMBERS PREVIOUS EXTREME POINTS AND SEARCHES IN THE CONVEX HULL OF THESE AND THE CURRENT ITERATE FOR THE NEXT ITERATE



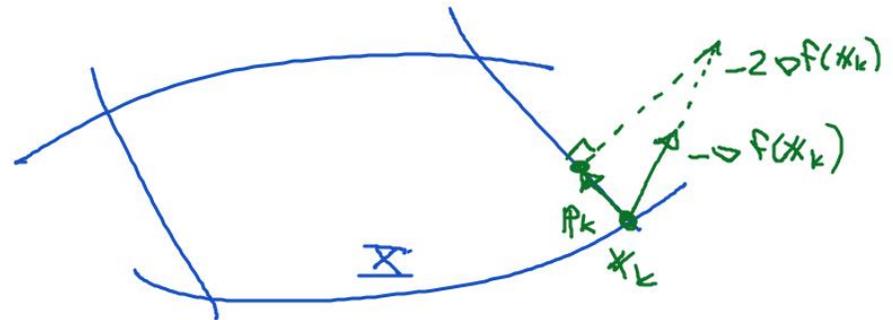
THE GRADIENT PROJECTION ALGORITHM

- IDEA: x^* LOCAL MIN $\Rightarrow x^* = \text{Proj}_{\mathcal{X}} [x^* - \nabla f(x^*)]$
- IF $x^* \neq \text{Proj}_{\mathcal{X}} [x^* - \nabla f(x^*)] \Rightarrow$ FEASIBLE DESC. DIR. CAN BE FOUND
- IF x_k IS NOT STATIONARY, THEN

$$p_k = \text{Proj}_{\mathcal{X}} [x_k - \alpha \nabla f(x_k)] - x_k \quad \text{FOR SOME } \alpha > 0$$

IS A FEASIBLE DESCENT DIRECTION

- IF $\mathcal{X} = \mathbb{R}^n$, THIS IS THE STEEPEST DESCENT METHOD



- HOW TO COMPUTE PROJECTIONS?

- ONLY POSSIBLE FOR SPECIFIC SETS

- HYPERCUBE $\{x \in \mathbb{R}^n \mid 0 \leq x_i \leq 1\}$

- UNIT SIMPLEX $\{x \in \mathbb{R}^n \mid \sum_{i=1}^n x_i = 1, x_i \geq 0\}$

- OTHER SPECIFIC TYPES OF SETS

- ANOTHER RELATED METHOD IS THE
GRADIENT PROJECTION METHOD WITH PROJECTION ARC

$$x_{k+1} = \text{Proj}_{\mathcal{X}} [x_k - \alpha_k \nabla f(x_k)]$$

- HOMEWORK: UNDERSTAND DIFFERENCE

CONVERGENCE

X IS NONEMPTY, CLOSED, CONVEX, AND $f \in C^1$ ON X .

IF x_k IS CHOSEN ACCORDING TO ARMIJOS RULE THEN

$\{x_k\}$ IS BOUNDED AND EVERY LIMIT POINT IS STATIONARY.

FRANK - WOLFE

X POLYHEDRON

GRADIENT PROJECTOR
METHOD

POSSIBLE TO PROJECT POINTS
ON X