## Examination, Finite automata and formal languages (DIT321/DIT322/TMV027/TMV028)

- Date and time: 2020-03-19, 14:00-18:00.
- Examiner: Nils Anders Danielsson. Telephone number: +46-31 7721680. Email address: nad@cse.gu.se.
- The GU grades Pass (G) and Pass with Distinction (VG) correspond to the Chalmers grades 3 and 5, respectively.
- To get grade $n$ on the exam you have to be awarded grade $n$ or higher on at least $n$ exercises.
- A completely correct solution of one exercise is awarded the grade 5. Solutions with minor mistakes might get the grade 5, and solutions with larger mistakes might get lower grades.
- Exercises can contain parts and/or requirements that are only required for a certain grade. To get grade $n$ on such an exercise you have to get grade $n$ or higher on every part marked with grade $n$ or lower (and every unmarked part), and you have to fulfil every requirement marked with grade $n$ or lower (as well as every unmarked requirement).
- Answers can be written in Swedish or English.
- Answers have to be saved to files in one of the following formats: PDF, JPEG or TXT. Use one file per question. Submit your solutions to Canvas before the deadline. Note that there is a separate Canvas assignment for each of the six questions. If Canvas is not working properly, send the solutions to the examiner using email, and include the course code in the subject header.
- Solutions can be rejected if they are hard to read (for instance if a picture is out of focus), unstructured, or poorly motivated.
- You do not need to provide proofs showing that algorithms covered in the course are correct. It is fine to use arguments of the following form: "Here I have used algorithm $X$ to compute the value $y$, and because the result of algorithm $X$ always satisfies property $P$, we have $P(y)$." However, you have to explain step by step why algorithm $X$ produces the value $y$.
- No collaboration is permitted, you have to work on your own.
- If you want to discuss the grading of the exam, contact the examiner no later than three weeks after the result has been reported.

1. Consider the context-free grammar $G=(\{S, A, B\},\{a, b\}, P, S)$, where the set of productions $P$ is defined in the following way:

$$
\begin{aligned}
& S \rightarrow A B \mid B S \\
& A \rightarrow a \mid A B \\
& B \rightarrow b
\end{aligned}
$$

(a) Is this grammar in Chomsky normal form?
(b) Construct the CYK table for $G$ and the string $a b b b$. No proof or explanation is required.
(c) For which nonterminals $N \in\{S, A, B\}$ do we have $a b b b \in L(G, N)$ ? Provide a proof.
2. Give a Turing machine $M$ such that $L(M)=\left\{0(12)^{n} 3 \mid n \in \mathbb{N}\right\}$. The machine's input alphabet should be $\{0,1,2,3\}$.
Explain your answer in such a way that the person correcting your exam can easily see that the answer is correct.
3. Consider the $\varepsilon$-NFA $A$ given by the following transition table:

|  | $a$ | $b$ | $\varepsilon$ |
| ---: | :--- | :--- | :--- |
| $\rightarrow s_{0}$ | $\emptyset$ | $\emptyset$ | $\left\{s_{1}\right\}$ |
| $s_{1}$ | $\left\{s_{1}, s_{2}\right\}$ | $\left\{s_{0}\right\}$ | $\emptyset$ |
| $* s_{2}$ | $\emptyset$ | $\left\{s_{0}, s_{1}\right\}$ | $\left\{s_{1}\right\}$ |

(a) Construct a regular expression $e$ over the alphabet $\{a, b\}$ satisfying $L(e)=L(A)$.
Prove that your construction is correct.
(b) Construct a minimal DFA $B$ over the alphabet $\{a, b\}$ satisfying $L(B)=L(A)$.
Prove that your construction is correct.
(c) Give a precise description, using natural language, of the language $L(A)$. Explain your answer in such a way that the person correcting your exam can easily see that the description is correct.
4. Do the regular expressions

$$
e_{1}=((\varepsilon+b) a)^{*}\left(a b(a+b a)^{*}\right)^{*}
$$

and

$$
e_{2}=(a+b a+a b)^{*}\left(b(a+b a+a b)^{*}\right)^{*}
$$

denote the same language over the alphabet $\{a, b\}$ ? Provide a proof.
5. Consider $X$ and $Y$, inductively defined subsets of $\{a, b\}^{*}$ :

$$
\overline{a \in X} \quad \frac{w \in Y}{a w \in X} \quad \frac{w \in X}{w b \in Y}
$$

(a) For grade 3: Define a context-free grammar $G$ satisfying $L(G)=X$. Explain your answer in such a way that the person correcting your exam can easily see that the answer is correct. (The explanation could refer to parts (b) and (c) below.)
(b) For grade 4: Prove that $X \subseteq L(G)$.
(c) For grade 4: Prove that $L(G) \subseteq X$.

For grade four it suffices to complete one of the proofs. For grade five both proofs are required.
6. Consider the following languages over $\{a, b, c\}$ :
(a) $\left\{u v\left|u, v \in\{a, b\}^{*},|u| \geq 3, \nexists v_{1}, v_{2} \in\{a, b\}^{*} . v=v_{1} a b v_{2}\right\}\right.$.
(b) $\left\{w c w \mid w \in\{a, b, c\}^{*}, \exists u, v \in\{a, b\}^{*} . w=u c v\right\}$.

For each language, answer the following two questions:

- Is the language regular?
- Is the language context-free?

Provide proofs. You are allowed to make use of the following lemmas/facts without proving them:

- The two pumping lemmas covered in the course.
- The closure properties covered in the course.
- The fact that for any alphabet $\Sigma$ with at least two elements the language $\left\{w w \mid w \in \Sigma^{*}\right\}$ is not context-free.
- The fact that if $\Sigma$ is an alphabet and $L \subseteq \Sigma^{*}$ is context-free, then, for any string $w \in \Sigma^{*},\left\{v \in \Sigma^{*} \mid w v \in L\right\}$ and $\left\{v \in \Sigma^{*} \mid v w \in L\right\}$ are also context-free (see "Quotients of Context-Free Languages" by Ginsburg and Spanier).

